# Polymorphism, Recursive Data Types, and Trees

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#### **Polymorphic types**

```
Consider the good old length function:
```

What is the type of length?

```
It could be int list -> int, or char list -> int, or even
(int list) list -> int! So it has many different types!
```

length should really work regardless of the type of the elements

```
It has type 'a list -> int, where 'a is a type variable
```

```
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```

#### 'a list -> int is a polymorphic type

length : 'a list -> int means that length has any type we can
obtain by replacing 'a with some arbitrary type

#### Examples:

'a list -> int is the most general type of length

The type system of F# gives the most general type, unless you give an explicit type declaration

*Type inference* is used to find this type

Some other polymorphic list functions (and lists):

```
List.head : 'a list -> 'a
List.tail : 'a list -> 'a list
take : int -> 'a list -> 'a list
drop : int -> 'a list -> 'a list
(@) : 'a list -> 'a list -> 'a list
(::) : 'a -> 'a list -> 'a list
[] : 'a list
```

# A Restriction for Polymorphic Types

Some polymorphic expressions are not allowed

Due to some deep technical reasons

This is called the "value restriction"

Affects expressions that are not value expressions

A value expression can be evaluated no further. Some examples:

17 [] (2.3,[]) sqrt [1;2;3] failwith

Some expressions that are not value expressions (can be evaluated further):

17+33 [] @ [] sqrt 5.0 List.head [1;2;3] failwith "Error!"

# The Value Restriction

The value restriction states that *right-hand sides in let declarations that are not value expressions can not be polymorphic* 

Some examples:

```
let a = 17 + x
\\ OK, 17 + x is not a value expression but has type int
let b = []
\\ OK, [] has polymorphic type 'a list but is a value expression
let c = [] @ []
\\ Not OK, [] @ [] has polymorphic type and is not a value expression
let d = 3 :: ([] @ [])
\\ OK, 3 :: ([] @ []) has (non-polymorphic) type int list
```

# **Recursive Data Types**

So far, we have defined data types with a number of cases, each of fixed size

How do we define data types for data like lists, which can have an arbitrary number of elements?

By making the data type definition *recursive*:

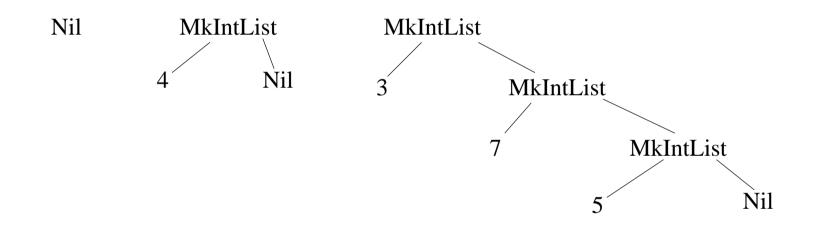
```
type IntList = Nil | MkIntList of (int * IntList)
```

An element of type IntList can be either Nil, or a data structure that contains an int and an IntList

Note similarity between data type declaration and context-free grammar

```
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```

Some IntList examples:



# Polymorphism

F#'s own data type for list is polymorphic

We can roll our own polymorphic list data type:

```
type List<'a> = Nil | MkList of 'a * List<'a>
```

Here, 'a is a type variable. Note the syntax < . . . > for user-defined polymorphic types: different from syntax for built-in polymorphic data types like 'a list

This data type is precisely the same as F#'s list data type, except that the constructor names are different!

Data type declarations can be recursive and polymorphic

Most of F#'s built-in data types can in principle be declared in the language itself

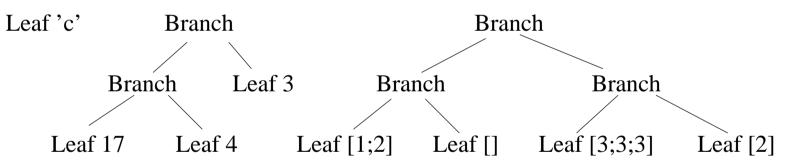
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## **Data Types for Trees**

We can easily make our own data types for *trees*, like:

type Tree<'a> = Leaf of 'a | Branch of Tree<'a> \* Tree<'a>

A data type for trees with data stored in the leaves



Many other variations possible, see examples in the book

Let us use this type for now

### **Operations on Trees**

Let us define some useful operations over our trees:

- a function to put the elements in a tree into a list,
- a function to compute the *size* (number of leaves) of a tree, and
- a function to compute the *height* of a tree.

(Code on next two slides)

To put the elements in a tree into a list:

Size (number of leaves):

#### Height:

### A Different Example: Arithmetic Expressions

Arithmetic expressions are really trees:

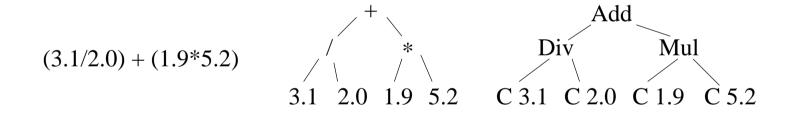
$$(3.1/2.0) + (1.9*5.2) \qquad \begin{array}{c} & & & \\ & &$$

Let us define a data type for arithmetic (floating-point) expressions! We can then use it for various symbolic manipulations of such expressions

(Data type declaration on next slide)

```
type Expr = C of float | Add of Expr * Expr | Sub of Expr * Expr
| Mul of Expr * Expr | Div of Expr * Expr
```

Each tree now represents an arithmetic expression:



## **Evaluating Expressions**

One operation is to evaluate expressions

```
eval : Expr -> float
let rec eval e =
  match e with
  | C x -> x
  | Add (e1,e2) -> eval e1 + eval e2
  | Sub (e1,e2) -> eval e1 - eval e2
  | Mul (e1,e2) -> eval e1 * eval e2
  | Div (e1,e2) -> eval e1 / eval e2
```

eval (Add ((C 17.0), Sub (C 3.0, C 1.0)))  $\implies$  19.0

eval is a simple interpreter for our expression trees

Exercise (mini-project): extend Expr with variables. Then define a small symbolic algebra package for manipulating and simplifying expressions, for instance:

- evaluate constant subexpressions
- simplify as far as possible using algebraic identities
- symbolic derivation
- etc...